# CMSC 332: Computer Networks Intro (cont.)

Professor Doug Szajda

### Chapter I: roadmap

- I. What is the Internet?
- I.2 Network edge
- I.3 Network core
- I.4 Delay & loss in packet-switched networks
- 1.5 Protocol layers and their service models
- **1.6** Networks Under Attack
- I.7 History of Computer Networking and the InternetI.8 Summary

### What Took You So Long?

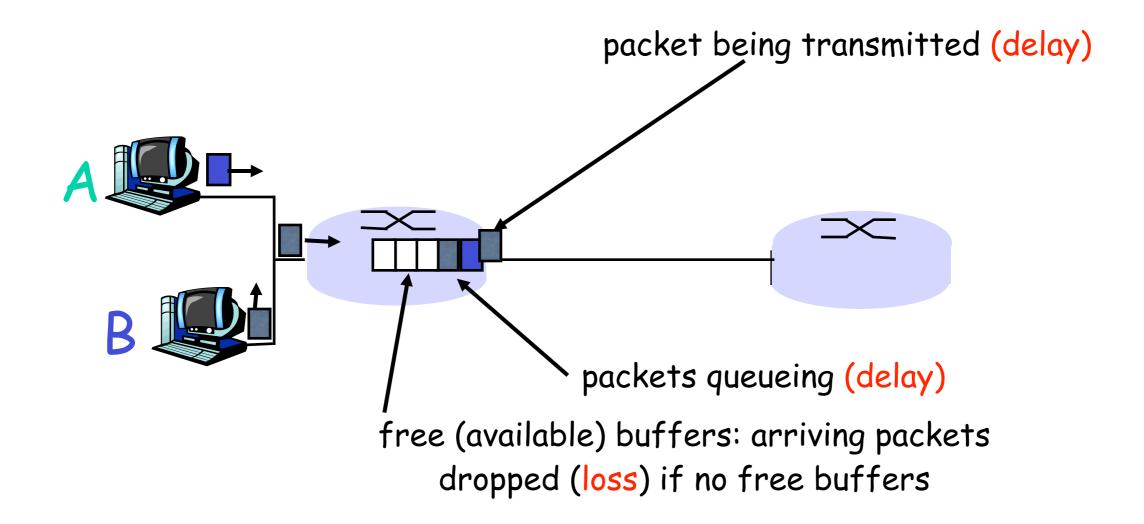
- The time it takes you to get to class depends on a lot of different factors.
  - How congested were the sidewalks? Any construction?
  - Was there a line outside the building? The classroom?
  - Were you carrying more things than usual?
- Network traffic is similarly influenced.
  - After all, traffic is not transmitted instantaneously.
  - Why?



# How do loss and delay occur?

packets queue in router buffers

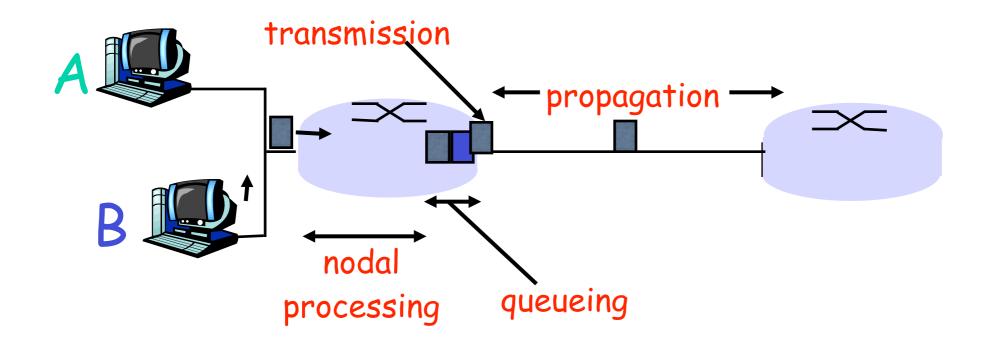
- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn



### Four sources of packet delay

- I. nodal processing:
  - check bit errors
  - determine output link

- 2. queueing:
  - time waiting at output link for transmission
  - depends on congestion level of router



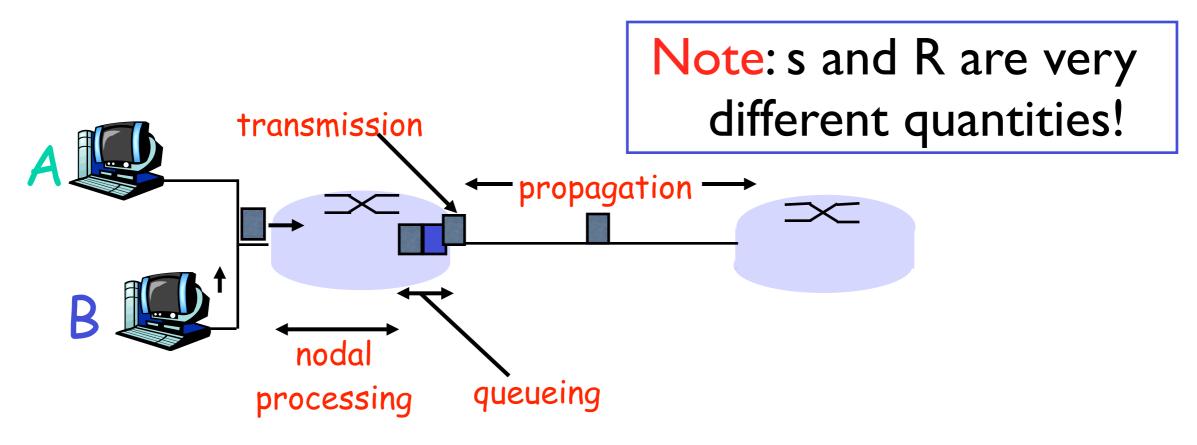
### Delay in packet-switched networks

#### 3. Transmission delay:

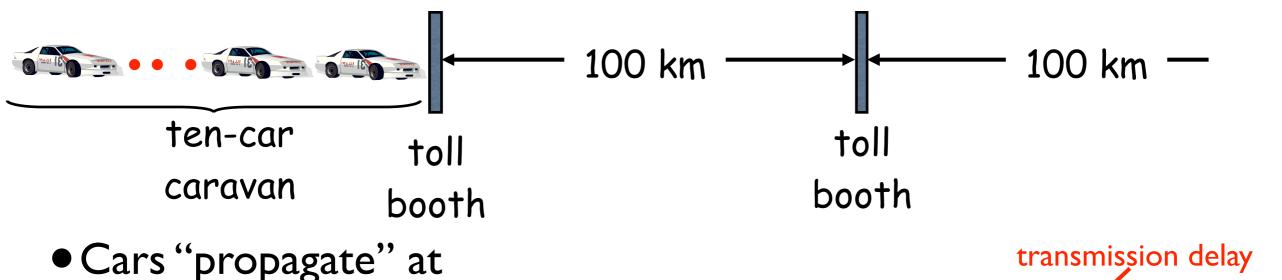
- R=link bandwidth (bps)
- L=packet length (bits)
- time to send bits into link = L/R

#### 4. Propagation delay:

- d = length of physical link
- s = propagation speed in medium (~2x10<sup>8</sup> m/sec)
- propagation delay = d/s



### Caravan analogy

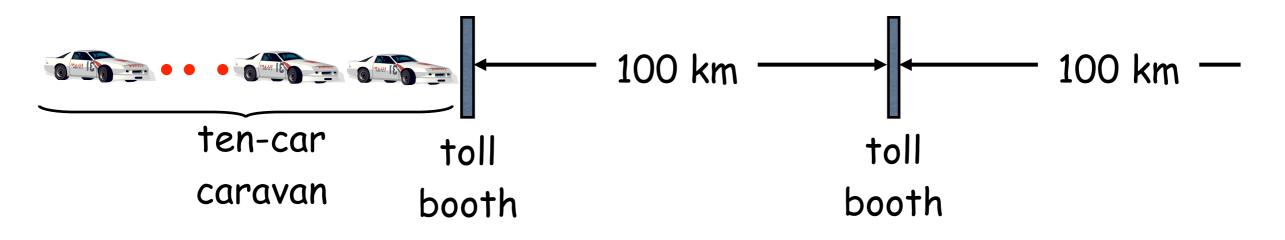


- Cars "propagate" at 100 km/hr
- Toll booth takes 12 sec to service a car (transmission time)
- car~bit; caravan ~ packet
- Q: How long until caravan is lined up before 2nd toll booth?

- Time to "push" entire caravan through toll booth onto highway = 12\*10 = 120 sec
- Time for last car to propagate from 1st to 2nd toll both: 100km/(100km/hr)= 1 hr
- A: 62 minutes

propagation delay

### Caravan analogy (more)



- Cars now "propagate" at 1000 km/hr
- Toll booth now takes 1 min to service a car
- Q:Will cars arrive to 2nd booth before all cars serviced at 1st booth?

- Yes! After 7 min, 1st car at 2nd booth and 3 cars still at 1st booth.
- Ist bit of packet can arrive at 2nd router before packet is fully transmitted at 1st router!
  - See Ethernet applet at the text book's Web site (K&R)

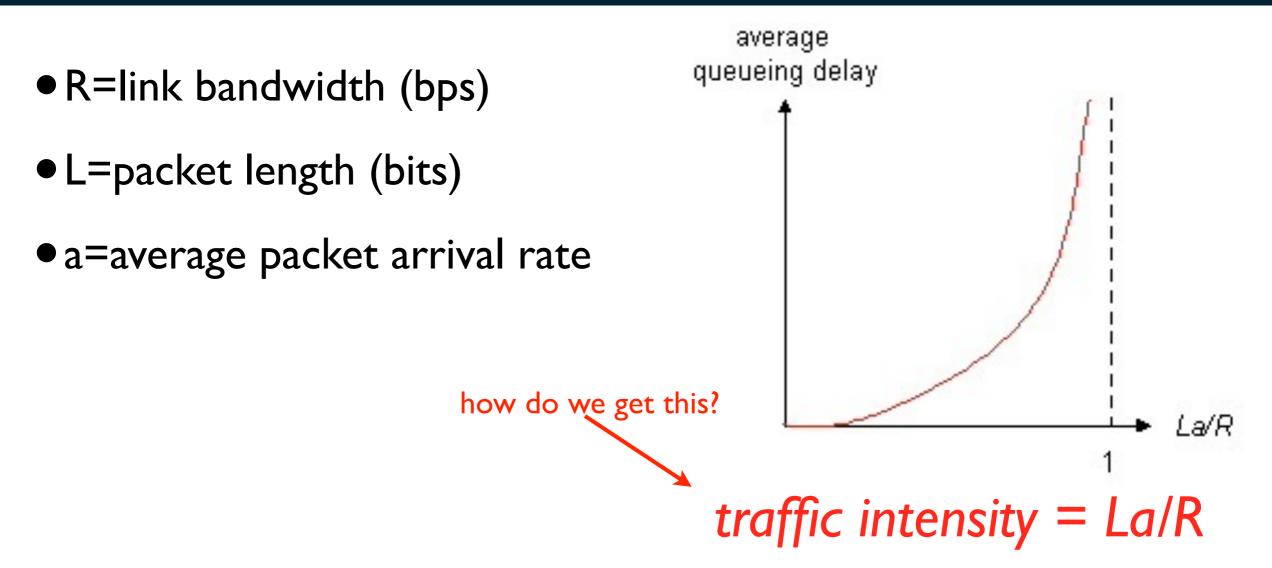
### Nodal delay

$$d_{\text{nodal}} = d_{\text{proc}} + d_{\text{queue}} + d_{\text{trans}} + d_{\text{prop}}$$

- $d_{proc}$  = processing delay
  - typically a few microsecs or less
- $d_{queue}$  = queuing delay
  - depends on congestion
- $d_{trans}$  = transmission delay
  - = L/R, significant for low-speed links
- $d_{prop}$  = propagation delay
  - a few microsecs to hundreds of msecs



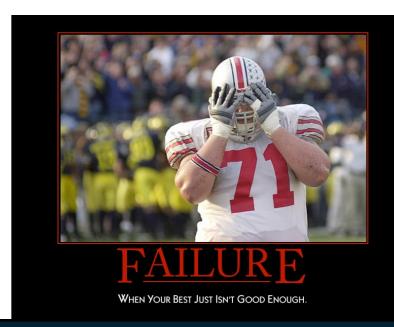
## Queueing delay (revisited)



- La/R ~ 0: average queueing delay small
- La/R -> I: delays become large
- La/R > I: more "work" arriving than can be serviced, average delay infinite!

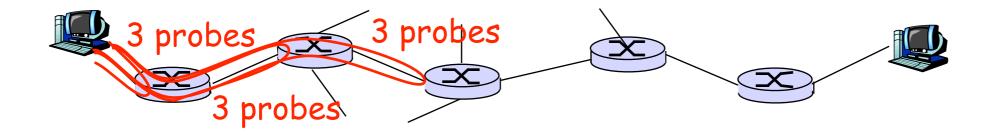
### Packet loss

- queue (aka buffer) preceding link has finite capacity
- when packet arrives to full queue, packet is dropped (aka lost)
- lost packet may be retransmitted by previous node, by source end system, or not retransmitted at all



### "Real" Internet delays and routes

- What do "real" Internet delay & loss look like?
- Traceroute program: provides delay measurement from source to router along end-end Internet path towards destination. For all *i*:
  - sends three packets that will reach router *i* on path towards destination
  - router *i* will return packets to sender
  - sender times interval between transmission and reply.



### "Real" Internet delays and routes

#### traceroute: gaia.cs.umass.edu to www.eurecom.fr

Three delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu 1 cs-gw (128.119.240.254) 1 ms 1 ms 2 ms 2 border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145) 1 ms 1 ms 2 ms 3 cht-vbns.gw.umass.edu (128.119.3.130) 6 ms 5 ms 5 ms 4 jn1-at1-0-0-19.wor.vbns.net (204.147.132.129) 16 ms 11 ms 13 ms 5 jn1-so7-0-0.wae.vbns.net (204.147.136.136) 21 ms 18 ms 18 ms 6 abilene-vbns.abilene.ucaid.edu (198.32.11.9) 22 ms 18 ms 22 ms trans-oceanic nycm-wash.abilene.ucaid.edu (198.32.8.46) 22 ms 22 ms 22 ms 8 62.40.103.253 (62.40.103.253) 104 ms 109 ms 106 ms link (how do we 9 de2-1.de1.de.geant.net (62.40.96.129) 109 ms 102 ms 104 ms 10 de.fr1.fr.geant.net (62.40.96.50) 113 ms 121 ms 114 ms know this?) 11 renater-gw.fr1.fr.geant.net (62.40.103.54) 112 ms 114 ms 112 ms 12 nio-n2.cssi.renater.fr (193.51.206.13) 111 ms 114 ms 116 ms 13 nice.cssi.renater.fr (195.220.98.102) 123 ms 125 ms 124 ms 14 r3t2-nice.cssi.renater.fr (195.220.98.110) 126 ms 126 ms 124 ms 15 eurecom-valbonne.r3t2.ft.net (193.48.50.54) 135 ms 128 ms 133 ms 16 194.214.211.25 (194.214.211.25) 126 ms 128 ms 126 ms 17 \* \* \* 18 \* \* \* \* means no response (probe lost, router not replying) 19 fantasia.eurecom.fr (193.55.113.142) 132 ms 128 ms 136 ms

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### Protocol "Layers"

Networks are complex!

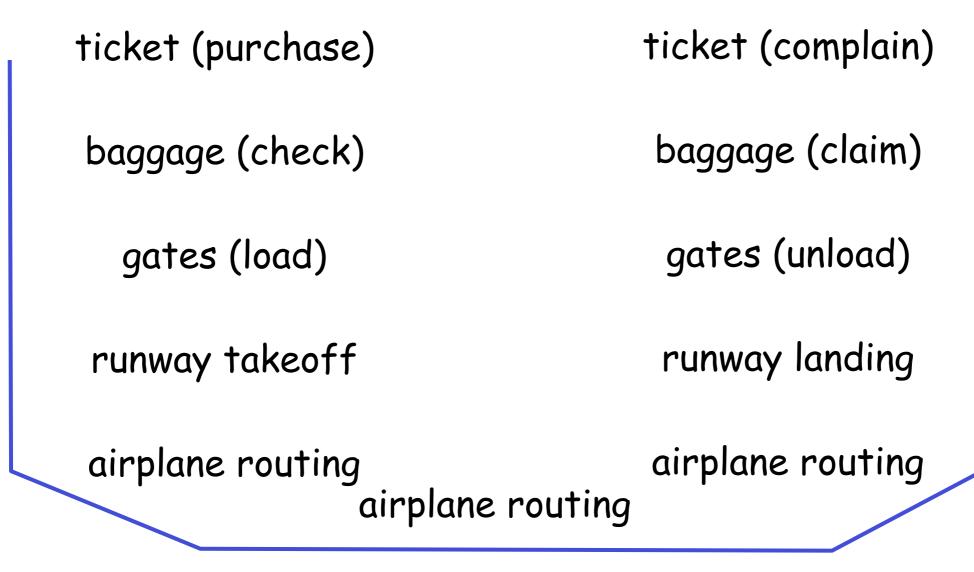
- many "pieces":
  - hosts
  - routers
  - links of various media
  - applications
  - protocols
  - hardware, software

#### Question:

Is there any hope of organizing structure of network?

Or at least our discussion of networks?

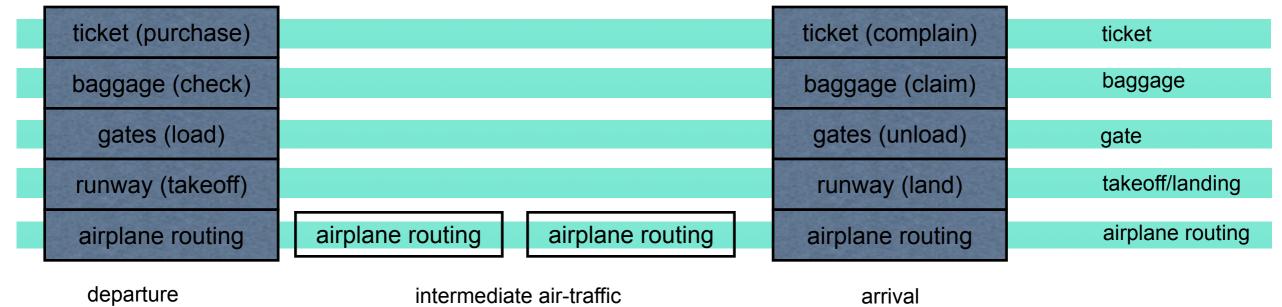
### Organization of air travel



• a series of steps

# Layering of airline functionality





airport

Layers: each layer implements a service

- via its own internal-layer actions
- relying on services provided by layer below

control centers

airport

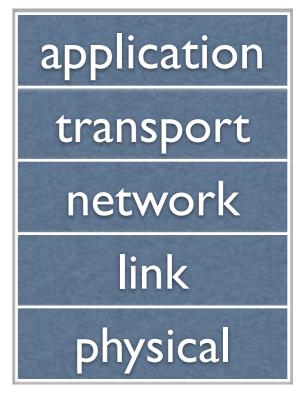
### Why layering?

Dealing with complex systems:

- explicit structure allows identification, relationship of complex system's pieces
  - layered reference model for discussion
- modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?

### Internet protocol stack

- application: supporting network applications
  - FTP, SMTP, HTTP
- transport: process-process data transfer
  - TCP, UDP
- network: routing of datagrams from source to destination
  - IP, routing protocols
- link: data transfer between neighboring network elements
  - PPP, Ethernet
- physical: bits "on the wire"

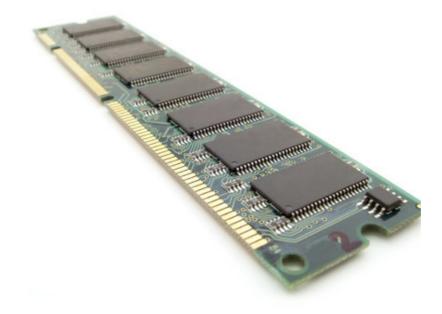


### **OSI** Reference Model

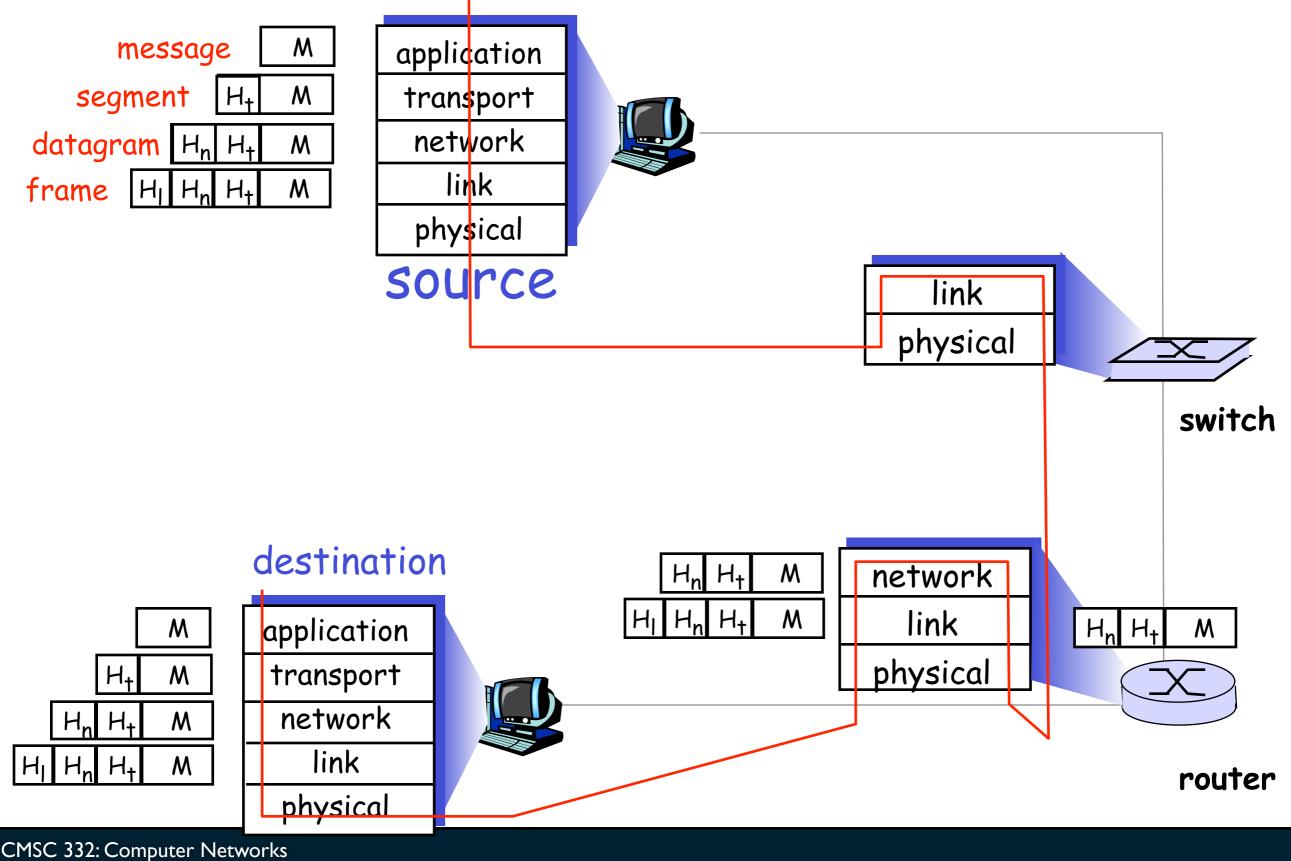
- The Open Systems Interconnection (OSI) model has two additional layers: Session and Presentation.
- Session Layer: Manages sessions between applications
  - (e.g., SSH, RTCP, RPC, NFS)
- Presentation Layer: Delivery and formatting of messages.
  - (e.g., RDP, ASCII)

### How Can I Remember This?

- There are a few simple mnemonics:
  - Please Do Not Tell Sales People Anything
  - All People Seem To Need Data Processing
  - Please Don't Nuke The South Pacific Again

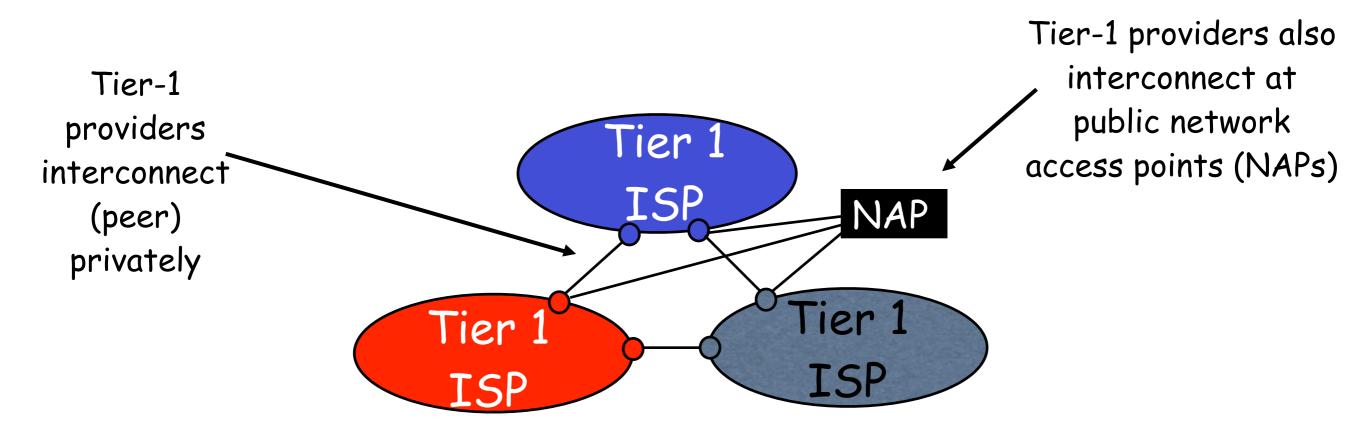


### Encapsulation

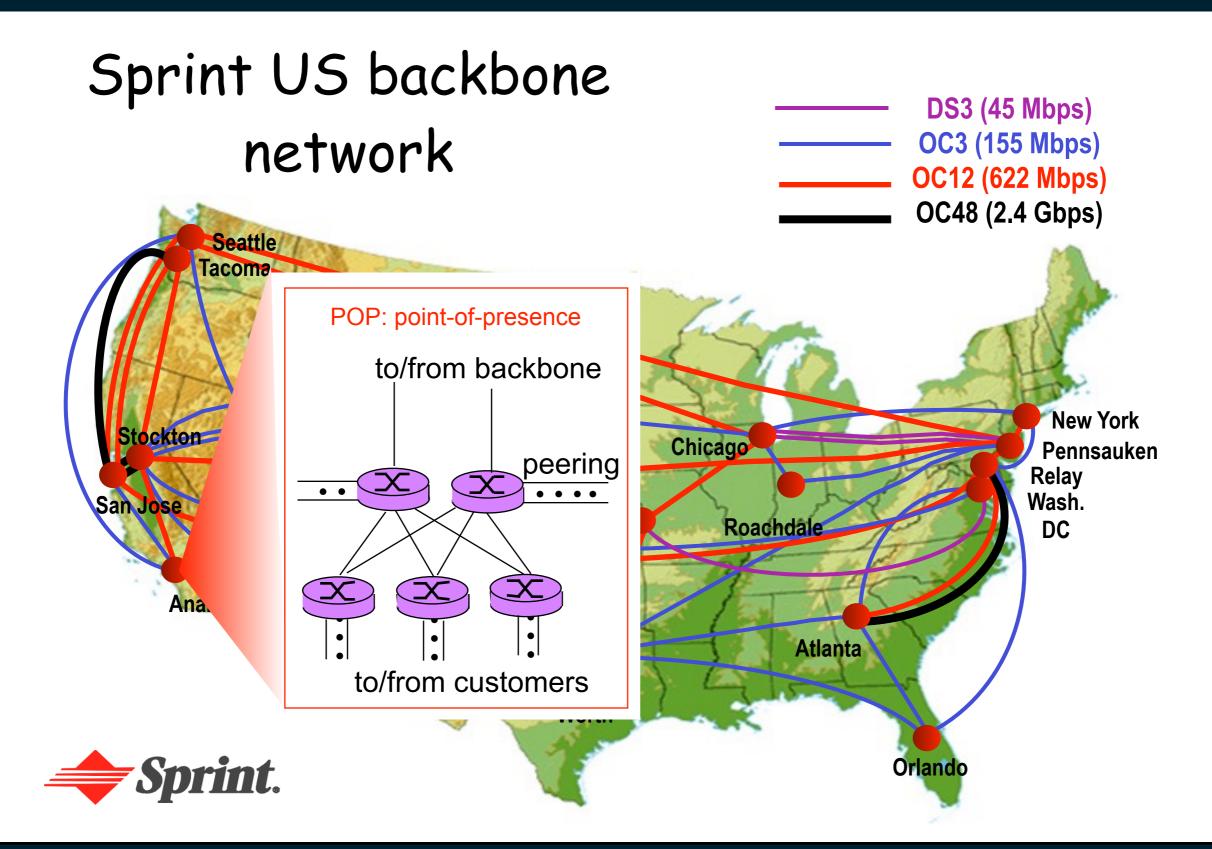


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- roughly hierarchical
- at center: "tier-I" ISPs (e.g., MCI, Sprint, AT&T, Cable and Wireless), national/international coverage
  - treat each other as equals

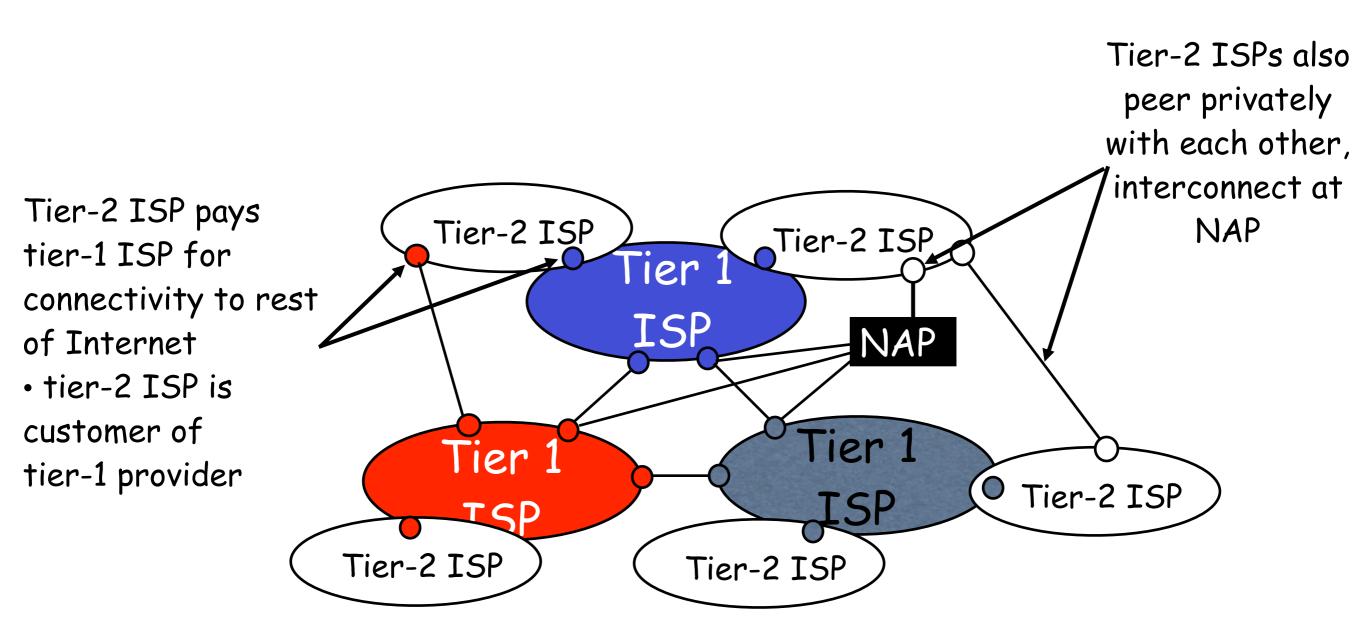


### Tier-I ISP: e.g., Sprint



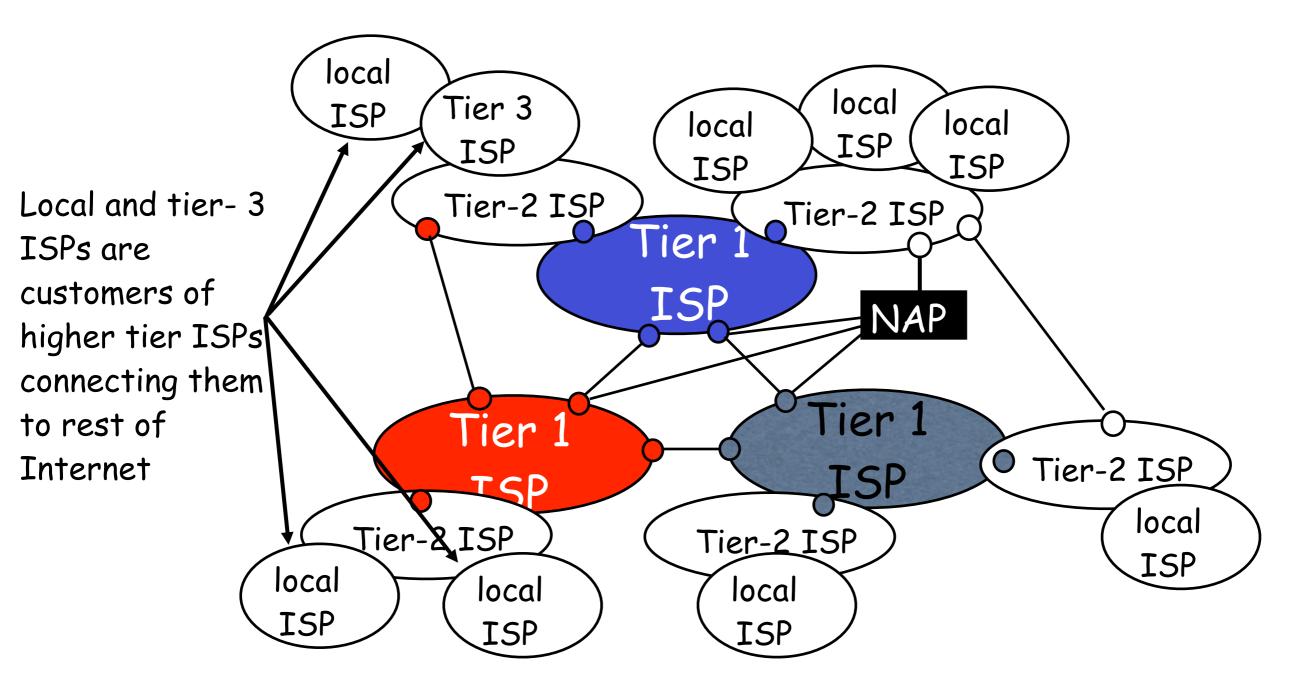
#### • "Tier-2" ISPs: smaller (often regional) ISPs

Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs

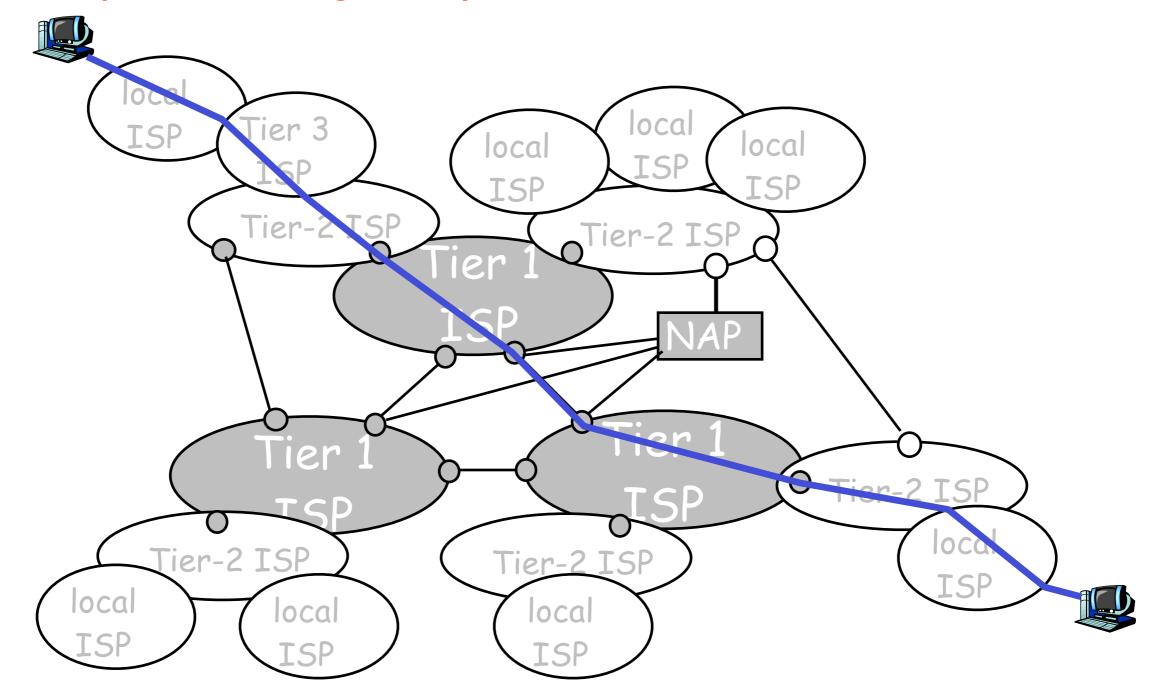


#### • "Tier-3" ISPs and local ISPs

last hop ("access") network (closest to end systems)



• a packet passes through many networks!



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### Back to the Silk Road

- When the Mongol Empire collapsed around 1400, most of the trade routes became too dangerous.
  - Columbus tried to avoid all of that.
- Any system in which huge amounts of money and information are exchanged will always be of interest to criminals.
  - The Internet is no different.
- What sorts of threats are out there?





- Malicious software is generically known as *malware*.
  - (e.g., Virus, worm, botnet, trojan horse)
- The distinction between these is often due to:
  - ...how they propagate...
  - ...what they control...
  - ...their usefulness to you...
- Anyone here every been infected?
- Anyone think they haven't?

### Attacking Availability

- An adversary may try to shut you down with a Denial of Service (DoS) attack.
- The book considers three categorizations, but the community has generally settled on two:
  - Flooding: Simply overwhelming your servers with more traffic than they can handle.
  - Logical: Exploiting a limited resource or known vulnerability.



### Packet Manipulation

- If the Internet is a network of networks, who says that someone in the middle can't mess with your packets?
- Assume that everything sent over the Internet is read or sniffed by someone/thing.
- Anyone with control of the wire can also arbitrarily drop or modify your packets.
  - When might this be a problem?
  - Does it happen?

### Authenticity

- How do you know who you are talking to?
  - In real life? On the web?
- Pretending to be someone else is easier than you think.
- An adversary can **spoof** identity in any number of ways...
  - Has this ever happened to you?



"On the Internet, nobody knows you're a dog."

### Security

- Security is, in general, a hard problem.
  - Even picking a good definition is difficult.
- As we move through this semester, ask yourself a few questions about the topics we study:
  - Did the designer consider malicious behavior?
  - How would I break this?
  - How would I fix this?



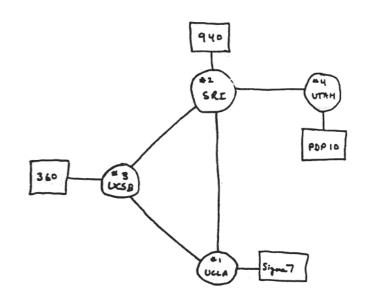
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- 1.4 Network access and physical media
- 1.5 Internet structure and ISPs
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- I.8 History

#### 1961-1972: Early packet-switching principles

- 1961: Kleinrock queueing theory shows effectiveness of packet-switching
- 1964: Baran packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational

- 1972:
  - ARPAnet public demonstration
  - NCP (Network Control Protocol) first host-host protocol
  - first e-mail program
  - ARPAnet has 15 nodes



THE ARPA NETWORK

#### 1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- late 70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:

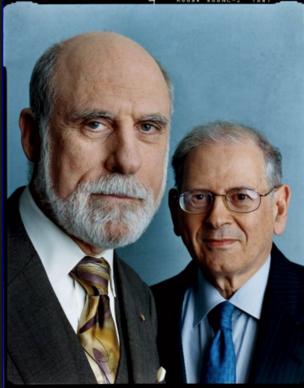
- minimalism, autonomy no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988:TCP congestion control

- new national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks



#### 1990, 2000's: commercialization, the Web, new apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s:Web
  - hypertext [Bush 1945, Nelson 1960's]
  - HTML, HTTP: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990's: commercialization of the Web

#### Late 1990's - 2000's:

- more killer apps: instant messaging, P2P file sharing, KaZaA
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

### Introduction: Summary

#### Covered a "ton" of material!

- Internet overview
- what's a protocol?
- network edge, core, access network
  - packet-switching versus circuit-switching
- Internet/ISP structure
- performance: loss, delay
- layering and service models
- history

#### You now have:

- context, overview, "feel" of networking
- more depth, detail to follow!



### Next Time

- Next Class
  - Read the "End-to-End argument" (link on website)
  - Read through Section 2.1
- Homework I will be assigned shortly

